

# Arena-Independent Finite-Memory Determinacy

## [Ongoing work]

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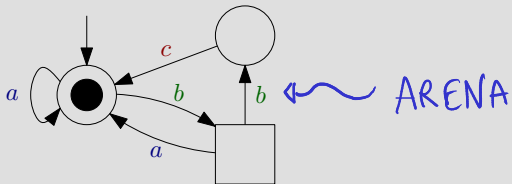
<sup>2</sup>F.R.S.-FNRS & UMONS – Université de Mons, Belgium

September 15, 2021 – Highlights 2021



# Games on graphs

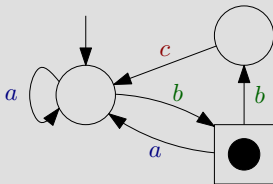
## Zero-sum turn-based games on graphs



# Games on graphs

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$$C = \{a, b, c\}$$

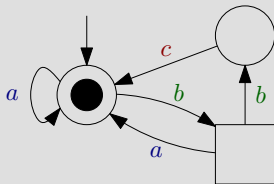


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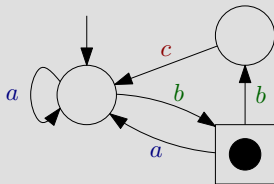


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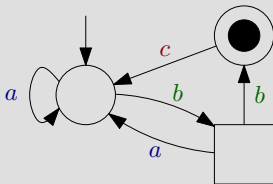
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$$C = \{a, b, c\}$$

$$W \subseteq C^{\omega}$$



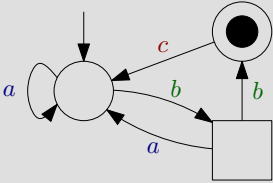
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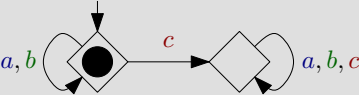
$$W \subseteq C^w$$



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## Strategy complexity

Given an objective, what information to remember about the past?



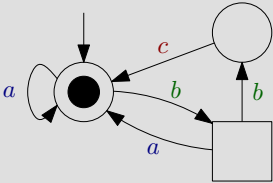
Memory structure  $\approx$  automaton.

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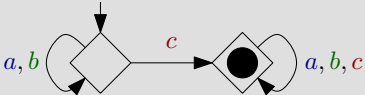
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babbcc...

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# Arena-independent finite-memory determinacy

## **Arena-independent finite-memory determinacy:**

objectives for which a **unique finite memory structure** suffices to play optimally in **all arenas**.

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objectives for which a **unique finite memory structure** suffices to play optimally in **all arenas**.

↪ good trade-off between **applicability** (e.g., all  $\omega$ -regular objectives, but not only) and **good properties**...

# One-to-two-player lifts

When does **two-player zero-sum** memory determinacy reduce to **one-player** memory determinacy?

Arenas \ Strat. comp.	Memoryless	Arena-ind. FM	Mildly growing
Finite deterministic	[GZ05] <sup>1</sup>	[BLORV20] <sup>2</sup>	[Koz21] <sup>3</sup>
Finite stochastic	[GZ09] <sup>4</sup>	[BORV21] <sup>5</sup>	
Infinite deterministic	[CN06, Kop08] <sup>6</sup>	<b>New work</b>	

<sup>1</sup>Gimbert and Zielonka, "Games Where You Can Play Optimally Without Any Memory", 2005.

<sup>2</sup>Bouyer, Le Roux, et al., "Games Where You Can Play Optimally with Arena-Independent Finite Memory", 2020.

<sup>3</sup>Kozachinskiy, "One-to-Two-Player Lifting for Mildly Growing Memory", 2021.

<sup>4</sup>Gimbert and Zielonka, "Pure and Stationary Optimal Strategies in Perfect-Information Stochastic Games with Global Preferences", 2009.

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# Infinite arenas, memoryless strategies

Let  $W \subseteq C^\omega$  be a **prefix-independent** objective.

## Theorem<sup>7</sup>

If **memoryless strategies** suffice to play optimally for both players in **one-player infinite deterministic arenas**, then  $W$  is a **parity condition**.

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In other words: there exists  $p: C \rightarrow \{0, \dots, n\}$  such that

$$w = c_1 c_2 \dots \in W \iff \limsup_i p(c_i) \text{ is even.}$$

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# Infinite arenas, arena-independent finite memory

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*using memory  $\mathcal{M}$*   
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*recognized by a parity automaton.*

Corollary 1. *New one-to-two-player lift.*

Corollary 2. *Strategic characterization of  $\omega$ -regular objectives.*

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Thanks! Questions?

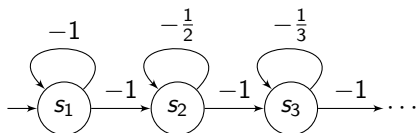


# Appendix

## Greater memory requirements in **infinite** arenas

Objective: get a mean payoff  $\geq 0$ .

- Memoryless strategies suffice in finite (even stochastic) arenas.
- Infinite memory is required in one-player deterministic infinite arenas.<sup>7</sup>



$\rightsquigarrow$  Possible to get 0 at the limit **with infinite memory**.

<sup>7</sup>Puterman, *Markov Decision Processes: Discrete Stochastic Dynamic Programming*, 1994.